

High Performance of MPI-FDTD on Multiprocessors

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Abstract: Parallel FDTD strategy is connected to analyze the electromagnetic issues of the electrically expansive targets on super computer. It is well known that the more the number of processors the less computing time expended. By the by, with the same number of processors, computing productivity is influenced by the plot of the MPI virtual topology. At that point, the impact of distinctive virtual topology plans on parallel execution of parallel FDTD is examined in detail. The common rules are displayed on how to get the most note worthy productivity of parallel FDTD calculation by optimizing MPI virtual topology. To appear the legitimacy of the displayed strategy, a few numerical comes about are given within the afterward portion.

Keywords: Parallel programs, MPI, FDTD and Computational Electromagnetics

Date of Submission: 04-11-2019

Date of acceptance: 22-11-2019

I. Introduction

(FDTD) by Yee, in 1966 [1], is one of the foremost well known three-dimensional strategies in computational electro-magnetics. The FDTD strategy has been connected to numerous calculations with exact comes about. Be that as it may, as a effective numerical procedure, the FDTD strategy is controlled to computation asset when analyzing the diffusing of electrically expansive targets. To overcome the issue, a Parallel FDTD calculation utilizing the Message Passing Interface (MPI) library was created by Volakis et al. in 2001 [2]. It is simple to actualize since the Yee conspire is express. The FDTD in Cartesian arranges can be effortlessly isolated into numerous subspaces, and each computer in a parallel framework bargains with one or a few sub-domains.

The FDTD calculation is combined with the MPI to run on parallel platforms. The MPI capacities are utilized to trade the digressive electric (attractive) fields on the boundary of the subdomain among the adjoining neighbors [3–5]. Parallel computation of the E-H components with an MPI Cartesian 2D topology is embraced and has been clarified step by step in [6]; as the creators of it tended to, it is the primary paper on Parallel FDTD utilizing MPI convention. Zhang et al. utilized an MPI Cartesian 3D topology in his inquire about [7]. As an expansion of these investigates, we have created a parallel FDTD code utilizing an MPI Cartesian 3D topology. It has been effectively run on PC cluster [5, 6] and edge server [8], both of which have a place to Science and Innovation on Recieving wire and Microwave Research facility in China.

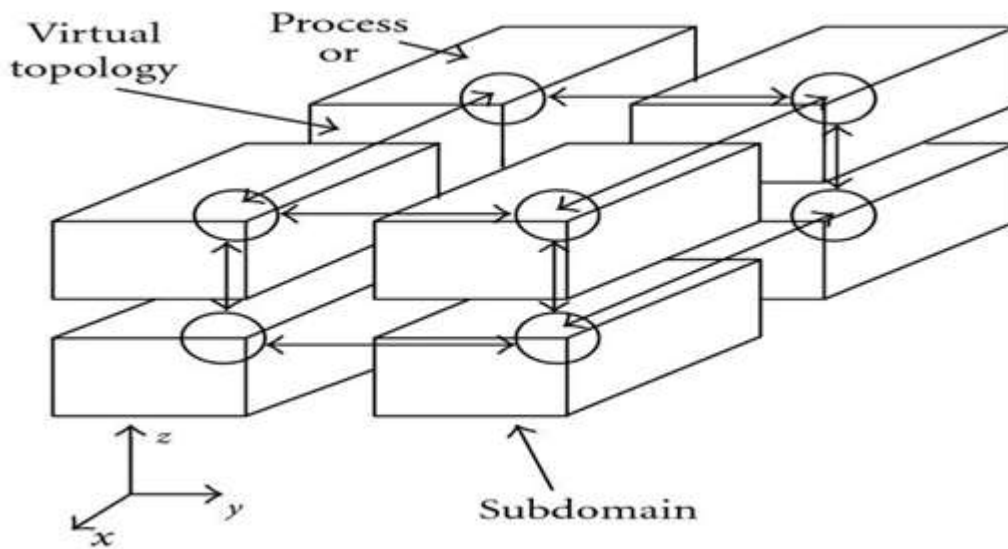
Presently, the code is created to unravel larger-scale electromagnetic issues. It is connected on super computer based on Linux framework, which has a place to Shanghai Supercomputer Center of China (SSC). Numerical illustrations demonstrate that the virtual topology will influence the computational proficiency of Parallel FDTD seriously. In this paper, the impact of diverse virtual topology plans on parallel execution of Parallel FDTD will be considered in detail. The common rules are displayed on how to get the most noteworthy proficiency of Parallel FDTD calculation by optimizing MPI virtual topology.

In Area 2, parallel FDTD calculation is displayed briefly. In Area 3, a few numerical comes about are given, which appear that this strategy is productive and exact.

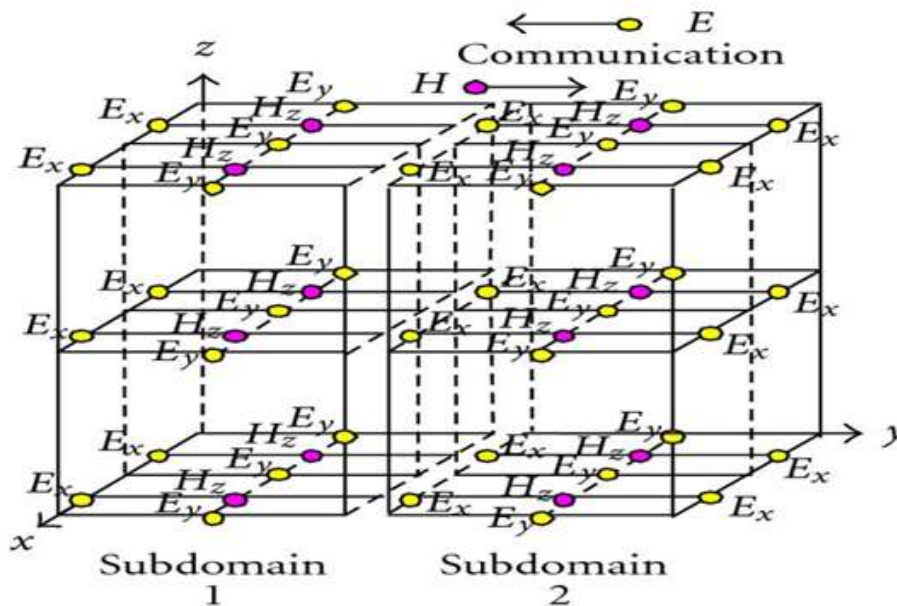
Current dispersion overthesurfaceof scramble andthe close field dispersion areplotted. Talks around the impact of distinctive virtual topology plans and diverse number of processors on parallel execution of Parallel FDTD are displayed in Segment 4. At long last, a few valuable conclusions are summarized. 2. Parallel Algorithm

MPI was proposed as a standard by a broadly based committee of sellers, implementers, and clients. Presently, it gets to be a definition for interfacing among a cluster of computers or the processors of a multiprocessor parallel computer. The key issue that MPI-based programming relates to is how to disperse the errands of clients to processors agreeing to the capability of each processor and decrease the communications among processors to as small as conceivable. Decreasing the communications is particularly vital as the speed of communication is distant slower than that of computation.

FDTD is simple to actualize since the Yee conspire is express. Other than, it has the rule advantage that since the lattice is normal and orthogonal, electromagnetic field components are effortlessly ordered by in Cartesian facilitates. The parallelism of the FDTD calculation is based on a spatial deterioration of the recreation issue geometry into coterminous nonoverlapping rectangular subdomains. The computational space can be effectively separated into about rise to parts along the three headings, and each processor in a parallel framework bargains with one or a few sub-domains, as appeared in Figure 1(a). The virtual topology of the processors' dispersion is chosen in a comparative shape as the issue segment. Each subdomain is mapped to its related hub where all the field components having a place to this subdomain are computed. To upgrade field values lying on interfacing between sub-domains, it is fundamental to trade information between neighboring processors. An x-z cut of the computational volume at one of the interfacing between hubs within the y-direction is appeared in Figure (b). In this paper, we embrace 3D communication design, which is presented as takes after. Below is the **illustration for the division and communication of the parallel FDTD in 3-D**.



(a) Simulated Partition



(b) Communication between the field components

A standard FDTD-algorithm is associated with our MPI to execute on the parallel systems. Important MPI subroutines and functions are applied for exchanging electric and magnetic fields on the subdomain boundaries and distributed among the axis of all fields.

Below is a step by step parallel algorithm and how it operates:

(1) Initialization.

- (a) MPI Initialization.
- (b) Reading the modeling parameters from the input files.
- (c) Creation of the three-dimensional Cartesian topology.
- (d) Creation of the inferred information sorts for communication.
- (e) Start time measurements.
- (f) Allocation of memory.
- (g) Setting all field components to zero.

(2) At each time step.

- (a) Exciting source as it were on processors that incorporate the source plane.
- (b) Calculation of modern attractive field components on each processor.
- (c) Communication of the attractive field components between processors.
- (d) Calculation of modern electric field components on each processor.
- (e) Communication of the electric field components between processors.
- (f) Calculation of transmission as it were on processors that incorporate the transmission plane.
- (g) Collection of field factors as it were on processors that incorporate location focuses.

(3) Reducing the transmission to a settled processor and composing it on file.

- (4) Saving comes about in files.
- (5) Deallocation of memory.
- (6) Stop time measurement.
- (7) MPI Finalization.
- (8) End.

3. Hardware Platform

(i) Think Station

Its machine-type model is 4155-D43 with a total of 24 Intel(R) Xeon(R) X5650 CPU cores (2.67 GHz per CPU) and a total of RAM approximately equal to 64 GB.

(ii) Shanghai Supercomputer Center (SSC)

The 37 nodes from Magic-cube Machine with a total of 512 AMD CPU cores (1.9 GHz per CPU and 4 cores on each CPU) 16 CPU cores on each node and 4 GB RAM per core, and a total of RAM approximately equal to 2.3TB. Infiniband is used for the network interconnection.

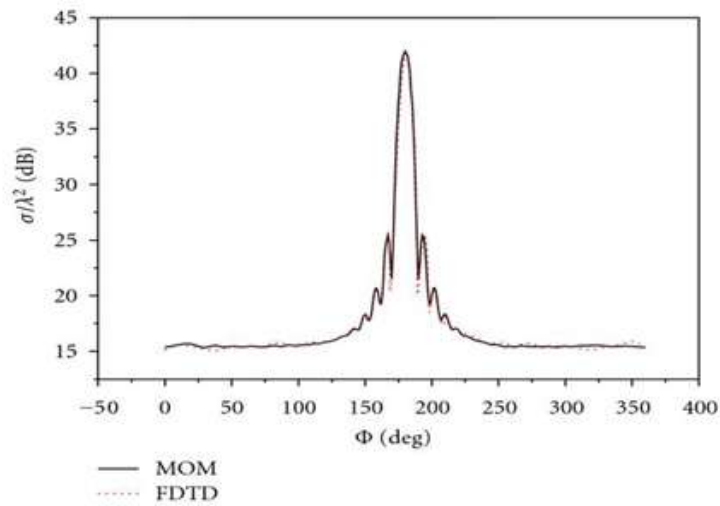
4. Numerical Results For the retaining UPML medium, we utilize a thickness of 5 cells within the taking after examples.

4.1. An Illustration for Validation For approval, the bistatic RCS is calculated for a PEC circle with a span, and the occurrence is arriving from the $-x$ axis and the polarization is along $-z$ axis. Working recurrence is 1.0 GHz. The increment

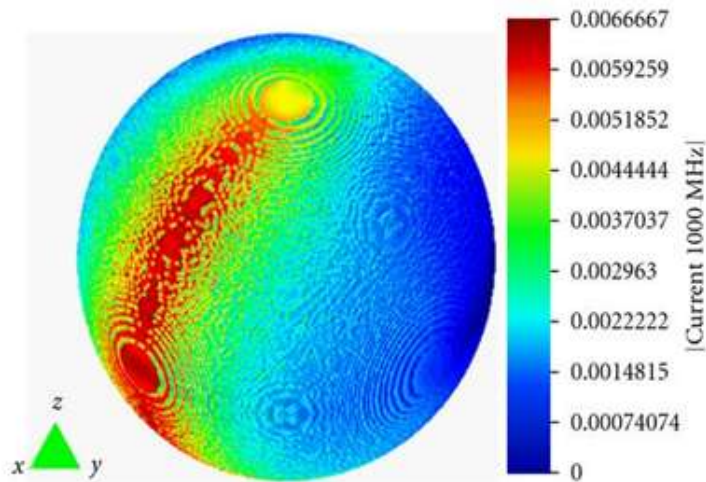
$dx = dy = dz = 0.02 \text{ m } (\lambda/15)$ utilized here, and the sum of FDTD networks is $144 \times 144 \times 144$. The bistatic RCS of the circle is appeared in Figure 2(a). The result is compared with current dissemination over the surface of the circle is given in Figure 2(b). In Figures 3(a) and 3(b), smooth form fills of the sufficiency of E field and H field dispersions in recurrence space are plotted, individually. This issue is calculated on Think Station. Add up to computation time (in seconds) with diverse number of processors and distinctive virtual topology plans in 1000 time-steps are compared in the below Table 1.

Table 1: Comparisons of computation time.

Cores	Virtual topology ($x \times y \times z$)	Computation time (in seconds)
1	$1 \times 1 \times 1$	646.25
2	$2 \times 1 \times 1$	330.75
4	$4 \times 1 \times 1$	172.25
4	$2 \times 2 \times 1$	170.88
8	$8 \times 1 \times 1$	98.38
8	$4 \times 2 \times 1$	92.75
8	$2 \times 2 \times 2$	91.25
16	$8 \times 2 \times 1$	95.12
16	$4 \times 4 \times 1$	91.38
16	$4 \times 2 \times 2$	89.00
24	$8 \times 3 \times 1$	69.38
24	$4 \times 3 \times 2$	67.38



(a) Checking the RCS x o y plane and the results by standard FDTD and its MOM



(b) random distribution over the sphere surface.

Figure 2

In Table 1, virtual topology plans are depicted as (X x Y x Z) for all three communication designs. On the off chance that the esteem is 1 in a few heading, it suggests that there's no topology in this heading. For example, 1 x 4 x 1 implies there's no topology in x and y bearings, individually, in 1 x 2 x 2 in this way the virtual topology is actually in one measurement. Essentially, implies there's no topology in x axis, in this way the virtual topology is really in two dimensions. From Table 1, it is apparent that expanding the number of processors can bring us the decrease of the computation time quickly. But distinctive virtual topology plans will fetched distinctive computation time indeed in the event that the code is run with the same number of processors.

We examine the parallel execution of the Parallel FDTD utilizing the distinctive dimensional virtual topology with the same number of processors. In this case, the computing time is calculated with the most limited time as the reference. Take the case of eight processors as the case. The reference calculating time which precisely is 91.25s with the 2 x 2 x 2 based on the comparisons:

$$8 \times 1 \times 1 (98.38 - 91.25) = 6.13s \text{ and } 4 \times 2 \times 1 (92.75 - 91.25) = 1.50s \dots \dots \dots (i)$$

From over, it is clear that for the same number of processors, the more the measurement of the virtual topology, the less the computation time required. Parallel effectiveness is appeared in Figure 4, in which the computing time expended by diverse number of processors is alluded as the most limited time for each case. Parallel proficiency is diminished with the expanding of processors. That's since the sums of the exchanged information are expanded with processors, and at that point the time expended on communication is expanded.

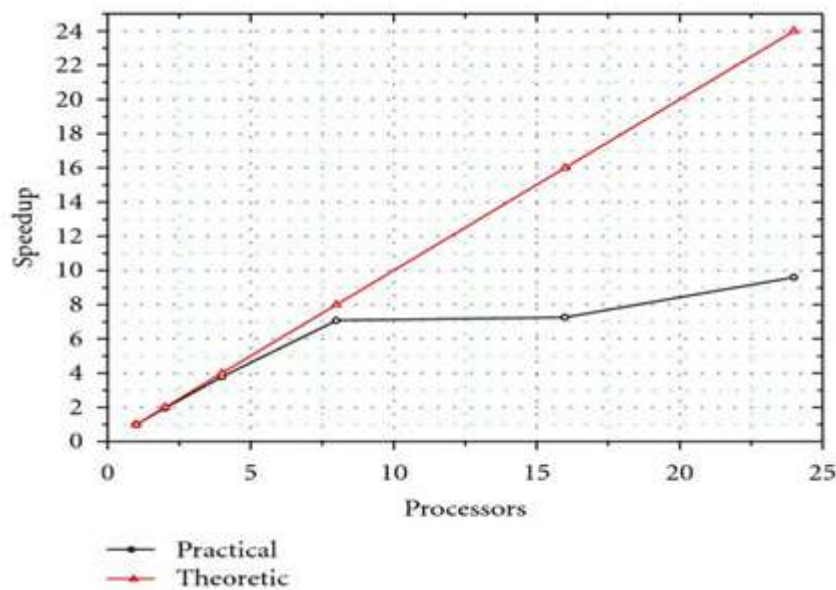


Figure 4: graphical efficiency of the process in parallel.

In expansion, the topology along the course where the sum of the FDTD networks is bigger can spare the computation time for the same dimensional virtual topology. Distinctive division subdomains with the same dimensional virtual topology lead to diverse sums of the exchanged information. Expression for the whole number of the networks lay on interface between processors are:

$$(nx-1) \times b \times c + (ny-1) \times a \times c + (nz-1) \times a \times b \dots \dots \dots (2)$$

Here a, b, c represents actual number of grids in x, y, z axis whereas nx, ny, nz are the values of integers and are subject to the instance:

$$nx \times ny \times nz = \text{processors, } nx \geq 1, ny \geq 1, nz \geq 1 \dots \dots \dots (3)$$

So, the topology conspire ought to be made along the headings where the sum of the FDTD networks is bigger to diminish the sums of the exchanged information, at that point to spare the communication time. Till presently, the common rules on how to get the most elevated effectiveness of Parallel FDTD calculation by optimizing MPI virtual topology can be drawn as follows. (1) If conceivable, the ideal virtual topology plot ought to be made in three measurements, and at that point the way better is in two measurements, which can bring us higher proficiency than in one dimension. (2) As to the same dimensional virtual topology, the topology conspire ought to be made along the bearings where the sum of the FDTD networks is bigger.

4.2. Radiation of the Waveguide with Ten Slots

A waveguide with ten openings is analyzed by parallel FDTD. The measurement of the waveguide and the space structure in this case are chosen as takes after: the thickness of the waveguide divider is 1.27 mm, the length of the space is 15.785 mm, the width of the opening is 2.54 mm, and all of the offsets of the spaces are 6.35 mm. Its FDTD demonstrate is appeared in Figure 5(a). Its working recurrence is 10 GHz. The change in $dx = dy = dz = 0.635\text{mm}$ is applied and the overall number of FDTD grids has dimension of an estimated $70 \times 50 \times 384$. The use of Gaussian pulse and using the update equations for the component in the phases:

$$E_y^{n+1}(i, j, k) = g(n\Delta t) + E_y^n(i, j, k) + \frac{\Delta t}{\epsilon_y(i, j, k)} \cdot \left[\frac{H_x^{n+1/2}(i, j, k+1) - H_x^{n+1/2}(i, j, k)}{\Delta z} - \frac{H_z^{n+1/2}(i+1, j, k) - H_z^{n+1/2}(i, j, k)}{\Delta x} \right], \dots\dots\dots (4)$$

$$g(n\Delta t) = \sin(\pi x/a) \sin(2\pi f_0 t) e^{-((t-t_0)/\sigma)^2}$$

We set to default $f_0 t_0$, and σ , ready to get a valuable recurrence transfer speed. At long last, smooth form fill of electric field dissemination in recurrence space on plane $y = 0.0$ Comes about of the radiation designs gotten by parallel FDTD. Add up to computation time (in seconds) with eight processors and distinctive virtual topology plans in 1000 time-steps is compared in Table 2.

Table 2: Comparisons of computation time.

Cores	Virtual topology ($x \times y \times z$)	Total grids lay on interface between processors	Computation time (in seconds)
8	$1 \times 1 \times 8$	24500	121.09
8	$1 \times 2 \times 4$	37380	124.56
8	$2 \times 1 \times 4$	29700	116.09
8	$2 \times 2 \times 2$	49580	118.00

With regards to table 2; the fastest execution time with 2-D is also brief and fast than when in 1-D. Therefore, the 3-D is not preferable than the 2-D. The converted data for $2 \times 2 \times 2$ scheme is greater than that of $2 \times 1 \times 4$. So we can conclude by saying the total computed average number of the grids in FDTD is the best and more reliable.

4.3. Investigation of the Scrambling of an Airplane

Then we analyze the scrambling of a superbly conducting plane whose FDTD show is appeared in Figure 7(a). Its working recurrence is 200 MHz. The change $dx = dy = dz = 0.08\text{m}$ is used. The heading of the occurrence wave is $-x$ direction and the polarization is along $+y$ direction. Inductive current conveyance over the plane surface is given in Figure 7(b). Figure 8 presents the bistatic RCS of the conducting plane gotten by utilizing Parallel FDTD. The recurrence space near-field conveyance on the plane $z = 0.05$ is given in Figure 9.

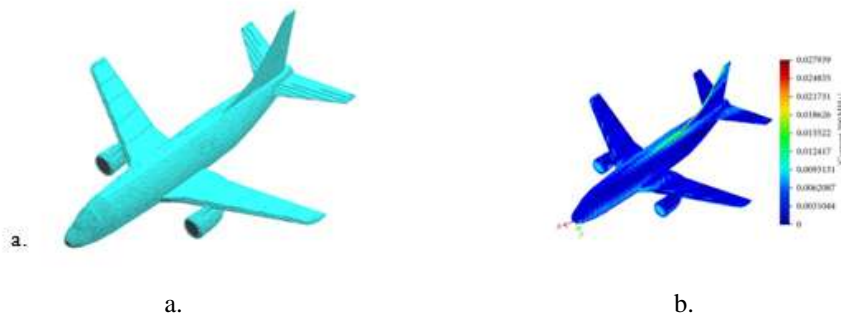


Figure 7

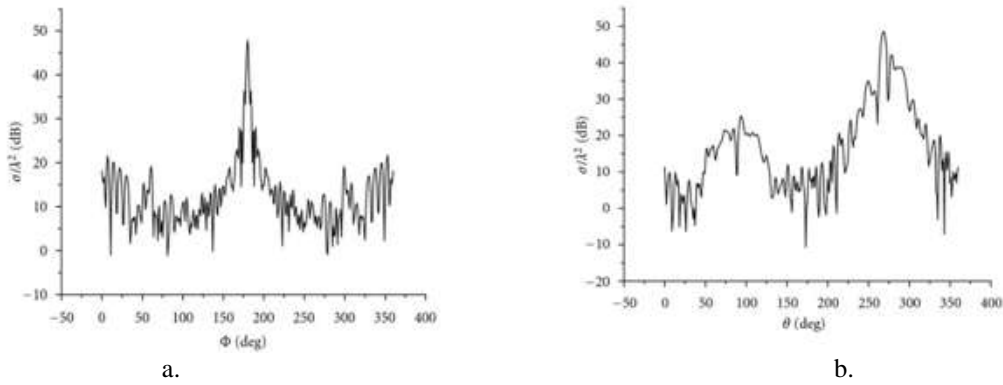
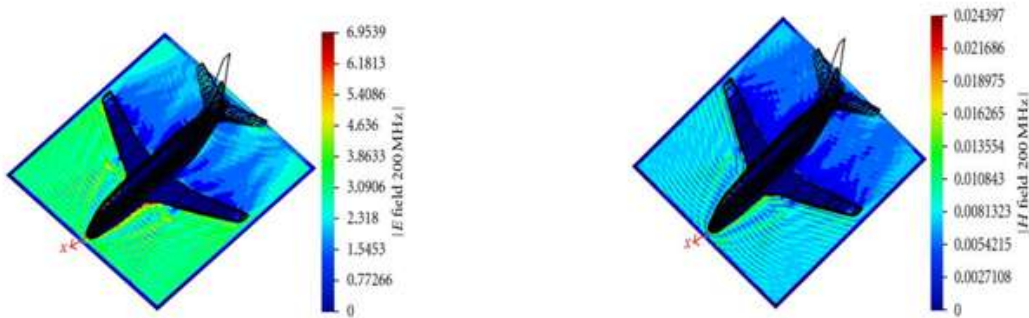


Figure 8



This illustration is calculated on super computer with 512 centers, which have a place with the same number of centers are recorded in Table 3. Add up to sum of the FDTD networks is 440 x 416 x 200. It is apparent that with the rules displayed some time recently. With the same dimensional virtual topology, the parallel execution is primarily influenced by the sums of the exchanged information particularly in large-scale issues.

Table 3: Comparisons of computation time.

Cores	Virtual topology (x × y × z)	Computation time (in seconds)	Total grids lay on interface between processors
512	8 × 8 × 8	896.00	2479680
512	8 × 16 × 4	3968.00	3000640

II. Conclusions

In this paper, parallel FDTD strategy is connected to analyze the electrically expansive targets. The code we created is effectively run on super computer in Shanghai Supercomputer Center of China (SSC). The impact of distinctive virtual topology plans on the parallel execution of Parallel FDTD is examined in profundity and in detail. The comes about appear that the computation time effectiveness can be moved forward by appropriately choosing MPI virtual topology plans. Taking after the two conclusions over, we will get the most noteworthy computational effectiveness.

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Adamu Abubakar "High Performance of MPI-FDTD on Multiprocessors" International Journal of Engineering Science Invention (IJESI), Vol. 08, No.11, 2019, PP 01-08